# A Fuzzy Approach To Three Stage Flow Shop Scheduling Model Including Job Block Criterion Using Branch And Bound Technique

K. Thangavelu<sup>1</sup>, G. Uthra<sup>2</sup>, S. Shunmugapriya<sup>3</sup>

P. G. and Research Department of Mathematics Pachaiyappa's College, Chennai-600 030, India. E-mail: kthangavelu14@gmail.com<sup>1</sup>, uthragopalsamy@yahoo.com<sup>2</sup>, priya010978@gmail.com<sup>3</sup>

#### **Abstract**

This paper deals with determining the optimal solution to three stage fuzzy flow shop scheduling problem. The processing times which are associated with probabilities are considered as triangular fuzzy numbers. The triangular fuzzy numbers are then defuzzified into crisp number using Yager's Ranking method. Branch and Bound technique is adopted to find the optimal sequence. The determination of optimal sequence by branch and bound method and minimum elapsed time is illustrated with a numerical example.

**Keywords:** Fuzzy flow shop scheduling problem, Yager's ranking method, Branch and Bound, Job block criteria.

#### Introduction:

In our day to day life, there are inadequate and vague information, since adequate data are not available in advance. In such circumstances it is better to provide approximate solutions with most and least possible values than with exact values. In such cases, fuzzy systems provide the better way to find proper solution.

Flow shop scheduling problem is one of classical problems widely used in manufacturing, production, management and so on. This problem aims at finding the optimal sequence of processing the jobs so as to minimize the total elapsed time. Johnson [10] proposed the well-known Johnson's rule in the two stage flow shop makespan scheduling problem.

Further the work was developed by Ignall and Scharge [6]. Brown and Lomnicki [2] studied the concept of flow shop scheduling using branch and bound method.

McCahon and Lee [12] analysed the job sequencing with fuzzy processing time. Ishibuchi and Lee [7] formulated fuzzy flow shop scheduling problem with fuzzy processing time. Hong and Chuang [5] introduced a triangular Johnson algorithm. Martin and Roberto [11] applied fuzzy scheduling to real time system. A valid approach based on statistical data for constructing a fuzzy flow shop sequencing is given by Yao and Lin [9]. TemizIzzettin and SerpilErol [8] introduced fuzzy branch and bound algorithm for flow shop scheduling.

Singh and Gupta [15], worked to minimize the rental cost in two stage flow shop problem. A new approach to two machine flow shop problem with uncertain processing time is given by Sanuja and Xueyan [14]. An heuristic approach to fuzzy flow shop problem on two machines is given by Singh, Sunita and Allawalia [16]. Sakthi et al [13] adopted Yager's ranking method to transform triangular fuzzy number to a crisp one. Deepak Gupta [3] applied branch and bound technique for three stage flow shop scheduling in which processing time and

set up time are both associated with probabilities including job block criteria.

D. Gupta, S. Sharma [4] studied specially structured two stage flow shop scheduling to minimize the rental cost. Ambika and Uthra [1] has worked with three stage flow shop scheduling problem using fuzzy processing times.

In this paper, triangular membership functions are used to represent fuzzy processing times associated with probabilities including job block criteria.

## 2. Definitions and preliminaries:

## 2. 1 Triangular fuzzy number:

A fuzzy number  $\tilde{a}$  on R is said to be a triangular fuzzy number or linear fuzzy number if its membership functions  $\tilde{a}$ :  $R \rightarrow [0, 1]$  has the following characteristics:

$$\widetilde{a}(x) = \begin{cases} \frac{x-a}{b-c} & \text{if } a \le x \le b \\ \frac{c-x}{c-b} & \text{if } b \le x \le c \end{cases}$$

$$0 & \text{otherwise}$$

The triangular fuzzy number is based on three value judgement: The minimum possible value a, the most possible value b and the maximum possible value c.

# 2. 2 α-cut and strong cut:

Given a fuzzy set A defined on X and any number  $\alpha \in [0, 1]$ , the  $\alpha$ -cut  $\alpha_A$ , and the strong  $\alpha$  cut  $\alpha_A^+$ , are the crisp sets,

$$\alpha_A = \{x/A(x) \ge \alpha\}$$
  
 $\alpha_A^+ = \{x/A(x) \ge \alpha\}$ 

# 2. 3 Yager's Ranking Formula

is defined by  $Y(\tilde{c}) = \int_0^1 0.5 (c_{\alpha}^L + c_{\alpha}^U)$ , where  $(c_{\alpha}^L, c_{\alpha}^U)$  is the  $\alpha$ -level cut of the fuzzy number  $\tilde{c}$ .

## 2. 4 Fuzzy arithmetic operations:

The following two operations are performed on triangular fuzzy numbers:

Let 
$$A=(a_1, a_2, a_3)$$
 and  $B=(b_1, b_2, b_3)$   
Addition:  $A+B=(a_1+b_1, a_2+b_2, a_3+b_3)$   
Subtraction:  $A-B=(a_1-b_3, a_2-b_2, a_3-b_1)$ 

## 2. 5 Notations:

 $a_i$ : Processing time for  $i^{th}$  job on machine A

**b**<sub>i</sub>: Processing time for *i*<sup>th</sup> job on machine B

 $c_i$ : Processing time for  $i^{th}$  jsob on machine C

 $p_i$ : Probability associated with processing time  $a_i$ 

 $q_i$ : Probability associated with processing time  $b_i$ 

 $r_i$ : Probability associated with processing time  $c_i$ 

 $A_i$ : Expected processing time for  $i^{th}$  job on machine A

 $\boldsymbol{B}_{i}$ : Expected processing time for  $i^{th}$  job on machine B

 $C_i$ : Expected processing time for  $i^{th}$  job on machine C

 $C_{ii}$ : Completion time for  $i^{th}$  job on machines A, B and C.

 $S_0$ : Optimal sequence

 $I_r$ : Partial schedule of r scheduled jobs.

I': The set of remaining (n-r) free jobs.

**\beta**: Equivalent job block.

# 2. 6 Assumptions:

O All the jobs are available for processing at time zero.

O Each job must be completed when started.

O To make job on a second machine, it must be completed on the first machine.

O Machines may be idle.

O Setup times are included in processing times along with their probabilities.

## 3. Mathematical development:

Consider n jobs say i= 1, 2, 3.... n are processed on three machines A, B & C in the order ABC. Each job i (i=1, 2, 3...n) has fuzzy processing time expressed by triangular fuzzy numbers  $a_i$ ,  $b_i$  &  $c_i$  on each machine respectively, assuming their respective probabilities  $\boldsymbol{p}_i$ ,  $\boldsymbol{q}_i$  and  $\boldsymbol{r}_i$  such that  $0 \leq \boldsymbol{p}_i \leq$ 

$$1, \sum p_i = 1, \ 0 \le q_i \le 1, \sum q_i = 1, 0 \le r_i \le 1, \sum r_i = 1.$$

Let an equivalent job  $\beta$  is defined as (k, m) where k, m are any jobs among the given n jobs such that k occurs before job m in the order of job block (k, m). The mathematical model of the problem in matrix form can be stated as:

| Jobs | Mach  | nine A | Machine B |       | Machine C      |       |
|------|-------|--------|-----------|-------|----------------|-------|
| i    | $a_i$ | $p_i$  | $b_i$     | $q_i$ | $c_{i}$        | $r_i$ |
| 1    | $a_1$ | $p_1$  | $b_1$     | $q_1$ | $c_1$          | $r_1$ |
| 2    | $a_2$ | $p_2$  | $b_2$     | $q_2$ | $c_2$          | $r_2$ |
| 3    | $a_3$ | $p_3$  | $b_3$     | $q_3$ | c <sub>3</sub> | $r_3$ |
| -    |       |        |           |       |                |       |
| •    |       |        |           |       |                |       |
| N    | $a_n$ | $p_n$  | $b_n$     | $q_n$ | $c_n$          | $r_n$ |

Our objective is to obtain the optimal schedule of all jobs which minimize the total elapsed time using branch and bound technique.

## 3. 1 Algorithm:

## Step 1:

Calculate expected processing time  $A_i$ ,  $B_i$ , and  $C_i$  on machines A, B and C respectively as follows:

(i) 
$$A_i = a_i \times p_i$$
 (ii)  $B_i = b_i \times q_i$ 

(iii) 
$$C_i = c_i \times r_i$$

## Step 2:

For triangular fuzzy numbers, using Yager's Ranking formula find expected processing times for machines  $A_i$ ,  $B_i$  and  $C_i$ 

## Step 3:

Find the equivalent processing time of job block  $\beta = (k, m)$  for machines A, B and C denoted by  $A_{\beta}$ ,  $B_{\beta}$  and  $C_{\beta}$ 

$$A_{\beta} = A_{k} + A_{m} - \min (A_{m}, B_{k})$$

$$B_{R} = (B_{1} + B_{2})/2$$

Where  $B_1 = B_k + B_m$  min  $(A_m, B_k)$  and

$$B_2 = B_k + B_{m} - \min(B_m, C_k)$$

$$C_{\mathcal{B}} = C_{\mathbf{k}} + C_{\mathbf{m}} - \min(B_{\mathbf{m}}, C_{\mathbf{k}})$$

#### Step 4:

Replace the processing time of jobs k and m by a single equivalent job  $\beta$  with processing time  $A_{\beta}$ ,  $B_{\beta}$  and  $C_{\beta}$ .

#### Step 5:

Calculate a lower bound for the 3 machine make span problem where  $A_i$ ,  $B_i$ ,  $C_i$  are the processing times of the  $i^{th}$  job on machines A, B and C using the formulae

(1) 
$$l_1 = t(J_r, 1) + \sum_{i \in J'_r} A_i + \min(B_i + C_i)$$

(2) 
$$l_2 = t(I_r, 2) + \sum_{i \in I_r'} B_i + \min(C_i)$$

(3) 
$$l_3 = t(J_r, 3) + \sum_{i \in J'_r} C_i$$

#### Step 6:

Calculate L= max  $\{l_1, l_2, l_3\}$  and evaluate L first for the n classes of permutations,

i. e. for these starting with 1, 2, 3... n respectively having labeled the appropriate vertices of the scheduling tree by these values.

## **Step 7:**

Now explore the vertex with lowest label. Evaluate L for the (n-1) subclasses starting with this vertex and again concentrate on the lowest label vertex. Proceeding until the end of the tree represented by two single permutations is reached, the total work duration is evaluated. Thus the optimal schedule of the jobs is obtained.

## Step 8:

Prepare in-out table for the optimal sequence obtained in step 3 and get the minimum total elapsed time.

#### 4. Numerical illustration.

Consider 6 jobs 3 machine flow shop problem where processing time of the jobs described by triangular fuzzy numbers in which jobs 2, 4 are to be processed as a group job (2, 4). Our objective is to obtain an optimal schedule using branch and bound technique and finding the total elapsed time.

Table 1: Processing time represented as triangular fuzzy numbers

| Jobs | Machine A    |       | Machine B    |       | Machine C    |       |
|------|--------------|-------|--------------|-------|--------------|-------|
|      | $a_i$        | $p_i$ | $b_i$        | $q_i$ | $c_{i}$      | $r_i$ |
| 1    | (1, 3, 7)    | 0. 2  | (5, 7, 9)    | 0. 2  | (18, 20, 22) | 0.2   |
| 2    | (3, 5, 8)    | 0.2   | (27, 29, 32) | 0.2   | (2, 4, 7)    | 0.2   |
| 3    | (27, 29, 31) | 0.3   | (1, 5, 6)    | 0.2   | (1, 3, 7)    | 0.1   |
| 4    | (1, 2, 5)    | 0.1   | (1, 3, 7)    | 0.1   | (1, 4, 6)    | 0.1   |
| 5    | (1, 4, 6)    | 0. 1  | (8, 11, 12)  | 0.2   | (1, 5, 6)    | 0.2   |
| 6    | (1, 5, 6)    | 0. 1  | (1, 3, 5)    | 0. 1  | (1, 5, 6)    | 0.2   |

Table 2: Processing times multiplied with probabilities

| Jobs | A                  | В                  | С               |
|------|--------------------|--------------------|-----------------|
| 1    | (0.2, 0.6, 1.4)    | (1, 1.4, 1.8)      | (3.6, 4.0, 4.4) |
| 2    | (0.6, 1.0, 1.6)    | (5.4, 5.8, 6.4)    | (0.4, 0.8, 1.4) |
| 3    | (5. 1, 8. 7, 9. 3) | (0.2, 1.0, 1.2)    | (0.1, 0.3, 0.7) |
| 4    | (0.1, 0.2, 0.5)    | (0.1, 0.3, 0.7)    | (0.1, 0.4, 0.6) |
| 5    | (0.1, 0.4, 0.6)    | (1. 6, 2. 2, 2. 4) | (0.2, 1.0, 1.2) |
| 6    | (0.1, 0.5, 0.6)    | (0.1, 0.3, 0.5)    | (0.2, 1.0, 1.2) |

Table 3: Defuzzification by Yager's ranking method

| Jobs | A      | В     | С      |
|------|--------|-------|--------|
| 1    | 0.8    | 1.4   | 4. 0   |
| 2    | 1. 05  | 5. 85 | 0. 85  |
| 3    | 7. 95  | 0. 85 | 0. 35  |
| 4    | 0. 25  | 0. 35 | 0. 375 |
| 5    | 0. 375 | 2. 1  | 0. 85  |
| 6    | 0. 425 | 0.3   | 0. 85  |

Table 4: Blocking the jobs 2, 4 into a single job  $\beta$  by job block

| Jobs | A      | В     | С     |
|------|--------|-------|-------|
| 1    | 0.8    | 1. 4  | 4     |
| β    | 1. 05  | 5. 9  | 0. 87 |
| 3    | 7. 95  | 0. 85 | 0. 35 |
| 5    | 0. 375 | 2. 1  | 0. 85 |
| 6    | 0. 425 | 0.3   | 0. 85 |

**Table 5: Calculation of Lower bounds** 

| $I_r$         | $Lb({\color{red}m{J}_{r}})$ |
|---------------|-----------------------------|
| 1             | 11. 75                      |
| β             | 13. 87                      |
| 3             | 19. 35                      |
| 5             | 11. 75                      |
| 6             | 11. 8                       |
| 1β            | 11. 75                      |
| 13            | 18. 75                      |
| 15            | 11. 75                      |
| 16            | 11. 8                       |
| 1 <b>β</b> 3  | 13. 9                       |
| 1 <b>β</b> 5  | 12. 25                      |
| 1 <b>β</b> 6  | 11.8                        |
| 1 <b>β</b> 63 | 14. 02                      |
| 1 <b>β</b> 65 | 11.8                        |

Table 6: The in-out table showing the minimum elapsed time

| Jobs | A      |        | В    |        | С      |        |
|------|--------|--------|------|--------|--------|--------|
|      | In     | Out    | In   | Out    | In     | Out    |
| 1    | 0      | 0.8    | 0.8  | 2. 2   | 2. 2   | 6. 2   |
| β    | 0.8    | 1. 85  | 2. 2 | 8. 1   | 8. 1   | 8. 98  |
| 6    | 1. 85  | 2. 275 | 8. 1 | 8.4    | 8. 98  | 9. 83  |
| 5    | 2. 275 | 2. 65  | 8.4  | 10. 5  | 10. 5  | 11. 35 |
| 3    | 2. 65  | 10.6   | 10.6 | 11. 45 | 11. 45 | 11.8   |

## **Solution:**

#### Step 1:

The processing times are represented as triangular fuzzy numbers with their probabilities as shown in table 1.

## Step 2:

The processing times are multiplied with their probabilities as shown in table 2.

(i) 
$$A_i = a_i \times p_i$$
 (ii)  $B_i = b_i \times q_i$ 

(iii) 
$$C_i = C_i \times r_i$$

# Step 3:

The processing times are defuzzified into crisp ones by using Yager's ranking formula

Y  $(\tilde{c}) = \int_0^1 0.5 (c_\alpha^L + c_\alpha^U)$ , where  $(c_\alpha^L, c_\alpha^U)$  is the  $\alpha$ -level cut of the fuzzy number  $\tilde{c}$  as shown in Table 3.

## Step 4:

Find an equivalent job  $\beta$  for the job block (2, 4) by calculating their processing times as follows.

$$A_{\beta} = A_k + A_m - \min(A_m, B_k)$$

$$B_{\beta} = (B_1 + B_2) / 2$$
  
where  $B_1 = B_k + B_m$ - min  $(A_m, B_k)$   
And  $B_2 = B_k + B_m$ - min  $(B_m, C_k)$   
 $C_{\beta} = C_k + C_m$ - min  $(B_m, C_k)$ 

The reduced problem is as shown in table 4.

# Step 5:

Calculation of lower bounds follows from steps 5, 6 and 7 of Algorithm 3. 1 and is shown in table 5.

## Step 6:

The in-out table is prepared for the optimal sequence showing the minimum elapsed time. The optimal Sequence is  $S_0$ :  $1\beta653$ . The minimum total elapsed time is 11. 8 and is shown in table 6.

#### 5. Conclusion.

In this paper, the processing times are considered as triangular fuzzy numbers. The fuzzy scheduling problem is converted into crisp one by Yager's ranking formula. The optimal sequence is determined using Branch and Bound techniques including job block criterion. Thus the scheduling problem for determining the minimum elapsed time in fuzzy environment is very effective under uncertainty.

## **REFERENCES**

- [1] G. Ambika and G. Uthra, Branch and Bound Technique in flow shop scheduling using fuzzy processing times, vol 8, No. 2, 2014, 37-42.
- [2] A. P. G. Brown and Z. A. Lominicki (1966), Some applications of the branch and bound algorithm to the machine scheduling problem, Operational Research Quarterly, 17 (1966) 173-182.
- [3] D. Gupta (2011), Application of branch and bound technique for n x 3 flow shop scheduling in which processing time associated with their respective probabilities, Mathematical Modeling and Theory, 2(1) (2011) 31-36.
- [4] D. Gupta, S. Sharma and S. Bala, (2012). Specially structured two stage flow shop scheduling to minimize the rental cost, International Journal of Emerging trends in Engineering and Development, vol 2, no. 1, pp. 206-215.
- [5] T. Hong, and T. Chuang, (1999). New triangular fuzzy Johnson algorithm, Computer and Industrial Engineering, 36(1), 179-200.
- [6] E. Ignall and L. Schrage(1965), Application of the branch and bound technique to some flow shop scheduling problem, Operations Research, 13(3) (1965) 400-412.
- [7] H. Ishibuchi, and K. H. Lee, (1996), Formulation of fuzzy flow shop scheduling with fuzzy processing time, In Proceeding of IEEE International Conference on Fuzzy system, 199-205.

- [8] T. Izzettin and S. Erol, Fuzzy branch and bound algorithm for flow shop scheduling, Journal of Intelligent Manufacturing, 15 (2004) 449-454.
- [9] Jing-Shing Yao and FrengTsc Lin, (2002), constructing a fuzzy flow shop sequencing model based on statistical data, International Journal of Appropriate Reasoning, 29(3), 215-234.
- [10] S. M. Johnson(1954), Optimal two and three stage production schedules with setup times include, Naval Research Logistics Quarterly, 1(1) (1954) 61-68.
- [11] L. Martin, and T. Roberto (2001). Fuzzy scheduling with application to real time system, Fuzzy sets and Systems, 121(3), 523-535.
- [12] S. McCahon and E. S. Lee, (1990). Job sequencing with fuzzy processing times. Computer and Mathematics with applications, 19(7), 31-41.
- [13] Sakthi Mukherjee and KajilaBasu, "Application of Fuzzy Ranking Method for solving Assignment Problems with Fuzzy costs", International Journal of Computational and Applied Mathematics, ISSN 1819-4966 volume 5 Number3(2010). Pp. 359-368.
- [14] P. Sanuja and S. Xueyan, A new approach to two machine flow shop problem with uncertain processing time. Optimization and Engineering 2006, 7(3), 329-343.
- [15] T. P. Singh, D. Gupta, (2005), Minimizing rental cost in two stage flow shop, the processing time associated with probabilities including job block Reflections de ERA, 1(2), 107-120.
- [16] T. P. Singh, Sunita and P. Allawalia, (2009). Fuzzy flow shop problem on two machines with single transport facility-An heuristic approach, Arya Bhatta journal of mathematics and informatics, 1(1-2), 38-46.