

## **Dominator Coloring Number of Regular Graphs**

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### **Abstract**

Graph coloring and domination in graphs have plenty of applications in computer, communication, biological and social network. The dominator coloring of the graph is obtained by merging the concept of graph node coloring and domination. It is defined as a proper coloring of nodes in which each node of the graph dominates all nodes of at least one color class. The smallest number of colors used for dominator coloring of nodes of a graph is called the dominator coloring number of the graph. In this paper the dominator coloring number of some regular graphs are obtained.

### **1. INTRODUCTION**

A network or graph consists of a collection of nodes and edges. Graph theory techniques enable the analysis of complex structures in networks; hence have plenty of applications in domains such as biological network, computer network, communication network, social network, etc.

A subset  $D_s$  of the node set of graph  $G$  which is such that each node  $V - D_s$  is adjacent to at least one member of  $D_s$  is called the dominating set [1]. The dominating set has application in computer and communication networks such as cell-phone network, mobile ad-hoc network, wireless sensor network etc. for routing and broadcasting the information between the nodes to the mobile devices. In biological network, the importance of dominating sets could be seen in the study of the human PPI network which holds details of the protein that operate and enable the biological processes within the cell [2, 3, 4].

A proper coloring of a graph  $G$  is a function  $f: V \rightarrow Z_+$  such that for  $v, w \in V$ ,  $f(v) \neq f(w)$  whenever  $v$  and  $w$  are adjacent nodes in  $G$ . Among all proper colorings of nodes of  $G$ , the smallest number of colors used in a node coloring is called the

chromatic number, denoted by  $\chi(G)$ . Graph coloring is an optimization problem which has many applications in circuit testing, assignment of mobile frequency, map coloring, allocation of register etc. Graph coloring techniques have applications in biological networks; in particular to PPI networks to help improve the quality of initial protein complexes. Thus these findings support to improve the existing protein complex detection methods [5].

As domination and coloring have application in interconnection networks, the dominator coloring or  $\chi_d$  coloring of a graph which is the amalgamation of node coloring and domination shall also be used in the study of behavior and properties of interconnection networks; in particular to mobile ad-hoc networks, vehicular adhoc network, air traffic network etc. The  $\chi_d$  coloring of a graph is defined as a proper coloring of nodes in a graph such that every node of  $G$  dominates at least a color class. The  $\chi_d$  coloring number denoted by  $\chi_d(G)$  is smallest number of colors used in the  $\chi_d$  coloring of  $G$  [7].

The notion of dominator coloring or  $\chi_d$  coloring was introduced by Hedetniemi et al[6]. Gera et.al [7-9] also studied this concept and has found the dominator coloring of various graphs and its association with chromatic number and domination number. The dominator coloring number in  $P_4$ - free graphs was given by Chellali and Maffray [10]. Merouane [11, 12] gave some bounds for planar and star-free graphs, exact values for split graphs and a characterization of trees. Arumugam et.al [13, 14] showed various bounds on dominator coloring, investigated about the algorithmic aspects of dominator coloring and gave a bound on Mycielskian of graph. The dominator coloring or  $\chi_d$  coloring of Prism graph, Quadrilateral Snake, Triangle Snake, Barbell graph, Degree splitting graph, m-splitting graph and m-shadow graph of path graph, Circular arc overlap graphs, closed Sun graph, closed Helm graph, generalized Flower Snark, Triangular belt, Alternate Triangular belt, Hajós graph, Trampoline graph, cyclic snake graphs, central graph, middle graph, total graph and line graphs of various graphs were discussed in various papers by various authors [15-27].

A graph in which every node has the same valency or degree is known as a regular graph. The study of regular graphs has applications in coding theory and computer science. Regular graphs are hard to be attacked since almost all nodes look the same i.e., they all have the same number of neighbors.

In this paper the dominator coloring number of some regular graphs (constructed using an algorithm) are obtained.

## 2. DOMINATOR COLORING NUMBER OF REGULAR GRAPHS

As a matter of fact, instead of trying to generate all  $k$ -regular graphs, we considered the  $k$ - regular graph constructed using the following algorithm [28]

### Algorithm 2.1:

Condition:  $0 < k < n$ .

First all the nodes are placed in a circle.

If  $n$  is odd:

For  $i$  in range  $(1, k)$ :

If  $i$  is odd:

Make edge for each node  $x$  steps away  $(i)$

If  $n$  is even:

If  $k$  is even:

Count Even numbers  $= \frac{n}{2}$

For  $i$  in range (Count Even numbers):

Make edge for each node  $x$  steps away  $(i + 1)$

If  $k$  is odd:

Count odd numbers  $= \frac{n-1}{2} + 1$

For  $i$  in range (Count odd numbers):

Make edge for each node  $x$  steps away  $\left(\frac{n}{2}\right) - i$

i.e., a regular graph is constructed as explained below

**Case 1:**  $n$  is odd and  $k$  must be even.

All the nodes are placed in a circle. Then the nodes are joined  $x$  positions away where  $x$  is all odd numbers between 1 to  $k$ .

**Example:**

$n = 9$	$k = 2$	Produce edge for each node $x$ steps away 1
$n = 9$	$k = 4$	Produce edge for each node $x$ steps away 1 Produce edge for each node $x$ steps away 3
$n = 9$	$k = 6$	Produce edge for each node $x$ steps away 1 Produce edge for each node $x$ steps away 3 Produce edge for each node $x$ steps away 5
$n = 9$	$k = 8$	Produce edge for each node $x$ steps away 1 Produce edge for each node $x$ steps away 3 Produce edge for each node $x$ steps away 5 Produce edge for each node $x$ steps away 7

**Case 2:  $n$  is even and  $k$  is even**

All the nodes are placed in a circle. Then the nodes are joined  $x$  positions away where  $x$  is a range of numbers from 1 to  $\frac{n}{2}$  and the range of these numbers is limited by the number of even numbers from 1 to  $k$  including  $k$ .

**Example:**

$n = 10$	$k = 2$	Produce edge for each node $x$ steps away 1
$n = 10$	$k = 4$	Produce edge for each node $x$ steps away 1 Produce edge for each node $x$ steps away 2
$n = 10$	$k = 6$	Produce edge for each node $x$ steps away 1 Produce edge for each node $x$ steps away 2 Produce edge for each node $x$ steps away 3
$n = 10$	$k = 8$	Produce edge for each node $x$ steps away 1 Produce edge for each node $x$ steps away 2 Produce edge for each node $x$ steps away 3 Produce edge for each node $x$ steps away 4

**Case 3:  $n$  is even and  $k$  is odd**

All the nodes are placed in a circle. Then the nodes are joined  $x$  positions away where  $x$  is a range of numbers from  $\frac{n}{2}$  to 1 and the range of these numbers is limited by the number of odd numbers from 1 to  $k$  including  $k$ .

**Example:**

$n = 10$	$k = 1$	Produce edge for each node $x$ steps away 5
$n = 10$	$k = 3$	Produce edge for each node $x$ steps away 5 Produce edge for each node $x$ steps away 4
$n = 10$	$k = 5$	Produce edge for each node $x$ steps away 5 Produce edge for each node $x$ steps away 4 Produce edge for each node $x$ steps away 3
$n = 10$	$k = 7$	Produce edge for each node $x$ steps away 5 Produce edge for each node $x$ steps away 4 Produce edge for each node $x$ steps away 3 Produce edge for each node $x$ steps away 2
$n = 10$	$k = 9$	Produce edge for each node $x$ steps away 5 Produce edge for each node $x$ steps away 4 Produce edge for each node $x$ steps away 3 Produce edge for each node $x$ steps away 2 Produce edge for each node $x$ steps away 1

**Proposition 2.2:** If  $C_n$  is a cycle on  $n$  nodes where  $n \geq 3$  then its dominator coloring

number is given by [8]  $\chi_d(C_n) = \begin{cases} \left\lceil \frac{n}{3} \right\rceil & \text{if } n = 4 \\ \left\lceil \frac{n}{3} \right\rceil + 1 & \text{if } n = 5 \\ \left\lceil \frac{n}{3} \right\rceil + 2 & \text{otherwise} \end{cases}$

**Theorem 2.3:** If  $G$  is a connected 2 regular graph on  $n$  nodes where  $n \geq 3$  then its

dominator coloring number is given by  $\chi_d(G) = \begin{cases} \left\lceil \frac{n}{3} \right\rceil & \text{if } n = 4 \\ \left\lceil \frac{n}{3} \right\rceil + 1 & \text{if } n = 5 \\ \left\lceil \frac{n}{3} \right\rceil + 2 & \text{otherwise} \end{cases}$

Proof:

Given  $G$  is a connected 2 regular graph on  $n$  nodes. Since  $G$  is connected and degree of every node is exactly 2, there exists an Eulerian cycle. Thus all connected 2 regular graphs with  $n$  nodes are isomorphic to  $n$  cycle. Hence by proposition 2.2 we

have  $\chi_d(G) = \begin{cases} \left\lceil \frac{n}{3} \right\rceil & \text{if } n = 4 \\ \left\lceil \frac{n}{3} \right\rceil + 1 & \text{if } n = 5 \\ \left\lceil \frac{n}{3} \right\rceil + 2 & \text{otherwise} \end{cases}$

**Theorem 2.4:** If  $G$  is a disconnected 2 regular graph on  $n$  nodes with  $k$  components then its dominator coloring number  $\chi_d(G)$  satisfies the relation  $\max(\chi_d(G_i)) + k - 1 \leq \chi_d(G) \leq \sum_{i=1}^k \chi_d(G_i)$ .

Proof:

Let  $G$  be a disconnected 2 regular graph on  $n$  nodes with  $k$  components  $G_1, G_2, \dots, G_k$ . Let  $\chi_d(G_i)$  be the dominator coloring number of each of the  $k$  components  $G_1, G_2, \dots, G_k$  of  $G$ . Since each component  $G_i$  of  $G$  is connected and 2 regular, it must contain some cycle  $C_k, k \geq 3$ . In other words a disconnected 2 regular graph  $G$  on  $n$  nodes is isomorphic to disjoint union of cycles. Then

$$\chi_d(G) \leq \sum_{i=1}^k \chi_d(G_i) \quad (1)$$

Let  $G_i$  be the component of  $G$  with maximum dominator coloring number  $\chi_d(G_i)$ . For every  $j \neq i$ , each  $G_j$  needs atleast a new color to dominate the entire color class of  $G_j$ . Hence the  $(k - 1)$  components of  $G$  need atleast  $(k - 1)$  colors to dominate its entire color class. Therefore

$$\chi_d(G) \geq \max(\chi_d(G_i)) + k - 1 \quad (2)$$

Combining equations (1) and (2), we have the result  $\max(\chi_d(G_i)) + k - 1 \leq \chi_d(G) \leq \sum_{i=1}^k \chi_d(G_i)$ .

**Theorem 2.5:** For any  $n \in \mathbb{Z}_+$ ,  $n$  is even,  $n \geq 4$ , if  $G$  is a  $n-2$  regular graph on  $n$  nodes, then its dominator coloring number is  $\chi_d = \frac{n}{2}$

Proof:

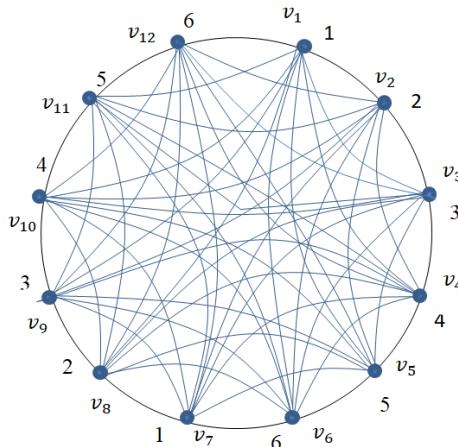
Let  $V(G) = \{v_i / 1 \leq i \leq n\}$  be the node set of the regular graph. The edges of the  $n-2$  regular graph are constructed using algorithm 2.1.

For dominator coloring, the nodes are assigned colors as follows:

The nodes  $v_i$  and  $v_{i+\frac{n}{2}}$  when  $1 \leq i \leq \frac{n}{2}$ , are given color  $i$ .

Then the nodes  $v_i$  and  $v_{i+\frac{n}{2}}$  when  $1 \leq i \leq \frac{n}{2} - 1$ , dominate color class  $i+1$ . The nodes  $v_n$  and  $v_{\frac{n}{2}}$  dominate color class 1.

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and the dominator coloring number of  $n-2$  regular graph on  $n$ (even) nodes is given by  $\chi_d = \frac{n}{2}$ .



**Figure 1:** Dominator coloring number of 10-regular graph on 12 vertices is  $\chi_d(G) = 6$

(Here  $n = 12$ ,  $n-2 = 10$  and  $\frac{n}{2} = 6$ )

**Theorem 2.6:** For any  $n \in \mathbb{Z}_+$ ,  $n$  is even,  $n \geq 4$ , if  $G$  is a  $\frac{n}{2}$  regular graph on  $n$  nodes then its dominator coloring number is  $\chi_d = \begin{cases} \frac{n}{2} & \text{when } n = 6 \text{ (or) } n \pmod 4 \equiv 0 \\ 4 & \text{when } n \neq 6 \text{ and } n \pmod 4 \equiv 2 \end{cases}$

Proof:

Let  $V(G) = \{v_i / 1 \leq i \leq n\}$  be the node set of the regular graph. The edges of the  $\frac{n}{2}$  regular graph are constructed using algorithm 2.1.

For dominator coloring, the nodes are assigned colors as follows:

$$\text{Let } q = \left\lceil \frac{n}{4} \right\rceil.$$

Case 1: When  $n = 6$

The nodes  $v_i, v_{i+q}$  and  $v_{i+2q}$  when  $1 \leq i \leq q$ , are given colors  $\left\lceil \frac{i}{q} \right\rceil, \left\lceil \frac{i}{q} \right\rceil + 1$  and  $\left\lceil \frac{i}{q} \right\rceil + 2$  respectively.

Then the nodes  $v_1, v_n$  dominate color class 2, the nodes  $v_2, v_3$  dominate color class 3 and the nodes  $v_4, v_5$  dominate color class 1 respectively.

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and  $\chi_d = \frac{n}{2}$  when  $n = 6$

Case 2: When  $n \pmod 4 \equiv 0$

The nodes  $v_i$  and  $v_{i+\frac{n}{2}}$  when  $1 \leq i \leq \frac{n}{2}$ , are given color  $i$ .

The nodes  $v_i$  and  $v_{i+\frac{n}{2}}$  when  $1 \leq i \leq \frac{n}{4}$ , dominate color class  $i + \frac{n}{4}$ . And the nodes  $v_i$  and  $v_{i+\frac{n}{2}}$  when  $\frac{n}{4} + 1 \leq i \leq \frac{n}{2}$ , dominate color class  $i - \frac{n}{4}$ .

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and  $\chi_d = \frac{n}{2}$  when  $n \pmod 4 \equiv 0$ .

Case 3: When  $n \pmod 4 \equiv 2$  and  $n \neq 6$

The nodes  $v_i, v_{i+q}$  and  $v_{i+2q}$  when  $1 \leq i \leq q$ , are given color  $\left\lceil \frac{i}{q} \right\rceil, \left\lceil \frac{i}{q} \right\rceil + 1$  and  $\left\lceil \frac{i}{q} \right\rceil + 2$  respectively. The nodes  $v_{n+1-i}$  when  $1 \leq i \leq n - 3q$ , are given color 4.

The node  $v_1$  dominates color class 2. The nodes  $v_{i+1}$  and  $v_{i+q+1}$  when  $1 \leq i \leq q$ , dominate color class  $\left\lceil \frac{i}{q} \right\rceil + 2$  and  $\left\lceil \frac{i}{q} \right\rceil + 3$  respectively. The nodes  $v_{n+1-i}$  when  $1 \leq i \leq q - 1$ , dominate color class 2. And the nodes  $v_{i+2q+1}$  when  $1 \leq i \leq q - 2$ , dominate color class 1.

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and  $\chi_d = 4$  when  $n \pmod 4 \equiv 2$  and  $n \neq 6$ .

Thus the dominator coloring number of  $\frac{n}{2}$  regular graph on  $n$ (even) nodes and  $n \geq 4$  is given by

$$\chi_d = \begin{cases} \frac{n}{2} & \text{if } n = 6 \text{ (or) } n \pmod 4 \equiv 0 \\ 4 & \text{if } n \pmod 4 \equiv 2 \text{ and } n \neq 6 \end{cases}.$$

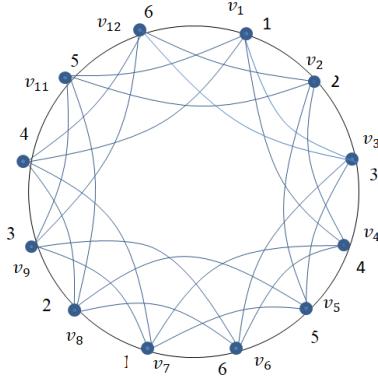


Figure 2(a)

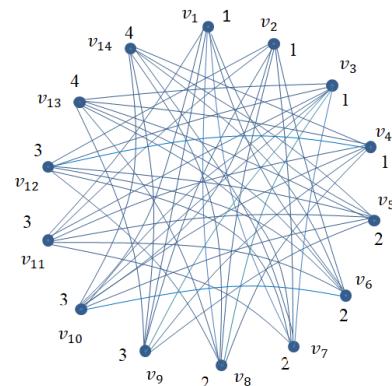


Figure 2(b)

**Figure 2(a):** Dominator coloring number of 6-regular graph on 12 vertices is  $\chi_d(G) = 6$   
(Here  $n = 12$  and  $n \pmod 4 \equiv 0$  )

**Figure 2(b):** Dominator coloring number of 7-regular graph on 14 vertices is  $\chi_d(G) = 4$   
(Here  $n = 14$  and  $n \pmod 4 \equiv 2$  )

**Theorem 2.7:** For any  $n \in \mathbb{Z}_+$ ,  $n$  is odd and  $n \geq 5$ , if  $G$  is a  $n - 3$  regular graph on  $n$  nodes, then its dominator coloring number is  $\chi_d = \left\lceil \frac{n}{2} \right\rceil$ .

Proof:

Let  $V(G) = \{v_i / 1 \leq i \leq n\}$  be the node set of the regular graph. The edges of the  $n - 3$  regular graph are constructed using algorithm 2.1.

For dominator coloring, the nodes are assigned colors as follows:

Let  $q = \left\lceil \frac{n}{4} \right\rceil$ .

Case 1: When  $n > 5$  and  $n \pmod 4 \equiv 3$

The nodes  $v_{4i-3}, v_{4i-1}$  when  $1 \leq i \leq q$ , are given color  $2i - 1$  respectively. The nodes  $v_{4i-2}, v_{4i}$  when  $1 \leq i \leq q$ , are given color  $2i$  respectively. The nodes  $v_{n-2}, v_n$  are given color  $\left\lceil \frac{n}{2} \right\rceil - 1$  and the node  $v_{n-1}$  is given color  $\left\lceil \frac{n}{2} \right\rceil$ .

The nodes  $v_{4i-3}, v_{4i-1}$  when  $1 \leq i \leq q$ , dominate color class  $2i$ . The nodes  $v_{4i-2}, v_{4i}$  when  $1 \leq i \leq q$ , dominate color class  $2i - 1$  respectively. The nodes  $v_{n-2}, v_{n-1}, v_n$  dominate color class  $\left\lceil \frac{n}{2} \right\rceil$ .

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and  $\chi_d = \left\lceil \frac{n}{2} \right\rceil$ .

Case 2: When  $n > 5$  and  $n \pmod 4 \equiv 1$

The nodes  $v_{4i-3}, v_{4i-1}$  when  $1 \leq i \leq q$ , are given color  $2i - 1$ . The nodes  $v_{4i-2}, v_{4i}$  when  $1 \leq i \leq q$ , are given color  $2i$  respectively. The node  $v_n$  is given color  $\left[ \begin{smallmatrix} n \\ 2 \end{smallmatrix} \right]$ .

The nodes  $v_{4i-3}, v_{4i-1}$  when  $1 \leq i \leq q$ , dominate color class  $2i$  and the nodes  $v_{4i-2}, v_{4i}$  when  $1 \leq i \leq q$ , dominate color class  $2i-1$  respectively. The node  $v_n$  dominates color class  $\left[\frac{n}{2}\right]$ .

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and  $\chi_d = \left\lceil \frac{n}{2} \right\rceil$ .

Case 3: When  $n = 5$

The nodes  $v_{4i-3}$ ,  $v_{4i-1}$  when  $1 \leq i \leq q$ , are given color  $2i - 1$ . The nodes  $v_{4i-2}$ ,  $v_{4i}$  when  $1 \leq i \leq q$ , are given color  $2i$  respectively. The node  $v_n$  is given color  $\left\lceil \frac{n}{2} \right\rceil$ .

The nodes  $v_{i+1}$  when  $1 \leq i \leq \left\lceil \frac{n}{2} \right\rceil$ , dominate color class  $i$ . And the nodes  $v_1$  and  $v_n$  dominate color class  $\left\lceil \frac{n}{2} \right\rceil$ .

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and and  $\chi_d = \left\lceil \frac{n}{2} \right\rceil$ .

Thus the dominator coloring number of  $n - 3$  regular graph on  $n$  (odd) nodes and  $n \geq 5$ , is given by  $\chi_d = \left\lceil \frac{n}{2} \right\rceil$ .

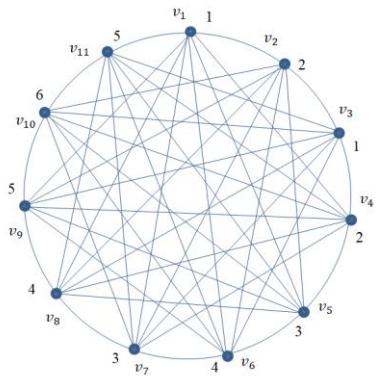


Figure 3(a)

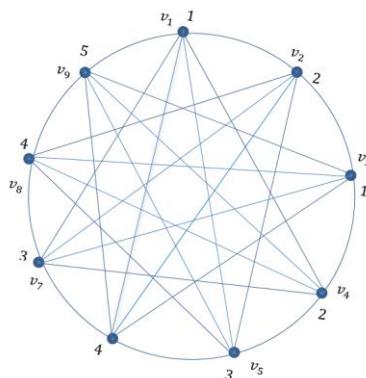


Figure 3(b)

**Figure 3(a):** Dominator coloring number of 8-regular graph on 11 vertices is  $\chi_d(G) = 6$

(Here  $n = 11$ ,  $n \pmod 4 \equiv 3$ ,  $n - 3 = 8$  and  $\chi_d(G) = \left\lceil \frac{n}{2} \right\rceil = 6$ )

**Figure 3(b):** Dominator coloring number of 6-regular graph on 9 vertices is  $\chi_d(G) = 5$

(Here  $n = 9, n \pmod 4 \equiv 1, n - 3 = 6$  and  $\chi_d(G) = \left[ \frac{n}{2} \right] = 5$ )

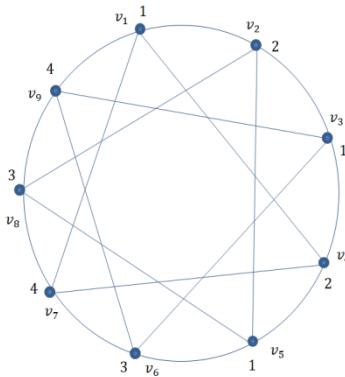
**Theorem 2.8:** For any  $n \in \mathbb{Z}_+$ ,  $n > 5$  and  $n \pmod{4} \equiv 1$ , if  $G$  is a  $\left\lfloor \frac{n}{2} \right\rfloor$  regular graph on  $n$  nodes, then its dominator coloring number is  $\chi_d = 4$ .

Proof:

Let  $V(G) = \{v_i / 1 \leq i \leq n\}$  be the node set of the regular graph. The edges of the  $\left\lfloor \frac{n}{2} \right\rfloor$  regular graph are constructed using algorithm 2.1.

For dominator coloring, the nodes are assigned colors as follows:

Let  $q = \left\lfloor \frac{n}{2} \right\rfloor$ . The nodes  $v_i$  when  $1 \leq i \leq q$ , are painted with color 1 when  $i$  is odd and color 2 when  $i$  is even. The node  $v_{\left\lfloor \frac{n}{2} \right\rfloor}$  is given color 1. The nodes  $v_{i+\left\lfloor \frac{n}{2} \right\rfloor}$  when  $1 \leq i \leq q$ , are painted with color 3 when  $i$  is odd and color 4 when  $i$  is even.



**Figure 4:** Dominator coloring number of 4 - regular graph on 9 vertices is  $\chi_d(G) = 4$ .

$$\left( \text{Here } n = 9, n \pmod{4} \equiv 1 \text{ and } \left\lfloor \frac{n}{2} \right\rfloor = 4 \right)$$

The nodes  $v_i$  when  $1 \leq i \leq q$ , dominate color class 2 when  $i$  is odd and color class 1 when  $i$  is even. The node  $v_{\left\lfloor \frac{n}{2} \right\rfloor}$  dominates color class 2. The nodes  $v_{i+\left\lfloor \frac{n}{2} \right\rfloor}$  when  $1 \leq i \leq q$ , dominate color class 4 when  $i$  is odd and color 3 when  $i$  is even.

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and the dominator coloring number of  $\left\lfloor \frac{n}{2} \right\rfloor$  regular graph on  $n$ (odd) nodes,  $n > 5$  and  $n \pmod{4} \equiv 1$  is given by  $\chi_d = 4$ .

**Theorem 2.9:** For any  $n \in \mathbb{Z}_+$ ,  $n \geq 5$  and  $n \pmod{4} \equiv 3$ , if  $G$  is a  $\left\lfloor \frac{n}{2} \right\rfloor$  regular graph on  $n$  nodes, then its dominator coloring number is  $\chi_d = 4$ .

Proof:

Let  $V(G) = \{v_i / 1 \leq i \leq n\}$  be the node set of the regular graph. The edges of the

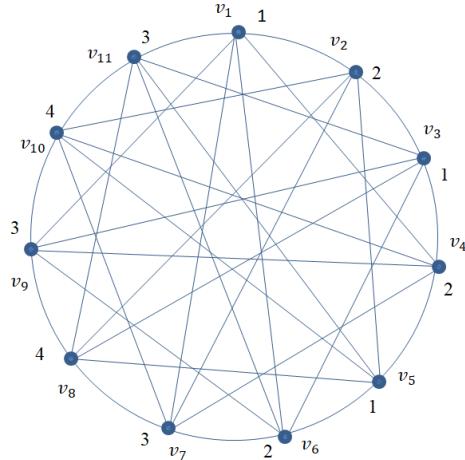
$\left[\frac{n}{2}\right]$  regular graph are constructed using algorithm 2.1.

For dominator coloring, the nodes are assigned colors as follows:

Let  $q = \left\lfloor \frac{n}{2} \right\rfloor$ . The nodes  $v_i$  when  $1 \leq i \leq q$ , are painted with color 1 when  $i$  is odd and color 2 when  $i$  is even. The nodes  $v_{i+q}$  when  $1 \leq i \leq q-1$ , are painted with color 3 when  $i$  is odd and color 4 when  $i$  is even.

The nodes  $v_i$  when  $1 \leq i \leq q$ , dominate color class 2 when  $i$  is odd and dominate color class 1 when  $i$  is even. The nodes  $v_{i+q}$  when  $1 \leq i \leq q-1$ , dominate color class 4 when  $i$  is odd and dominate color class 3 when  $i$  is even.

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and the dominator coloring number of  $\left[\frac{n}{2}\right]$  regular graph on  $n$ (odd) nodes,  $n \geq 5$  and  $n \pmod 4 \equiv 3$  is given by  $\chi_d = 4$ .



**Figure 5:** Dominator coloring number of 6 - regular graph on 11 vertices is  $\chi_d(G) = 4$ .

(Here  $n = 11$ ,  $n \pmod 4 \equiv 3$  and  $\left[\frac{n}{2}\right] = 6$ )

**Theorem 2.10:** For any  $n \in \mathbb{Z}_+$ ,  $n \geq 5$  and  $n \pmod 4 \equiv 1$ , if  $G$  is a  $\left[\frac{n}{2}\right] + 1$  regular graph on  $n$  nodes, then its dominator coloring number is  $\chi_d = 5$ .

Proof:

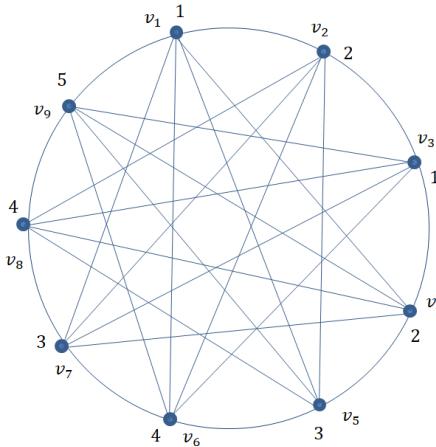
Let  $V(G) = \{v_i / 1 \leq i \leq n\}$  be the node set of the regular graph. The edges of the  $\left[\frac{n}{2}\right] + 1$  regular graph are constructed using algorithm 2.1.

For dominator coloring, the nodes are assigned colors as follows:

Let  $q = \left\lfloor \frac{n}{4} \right\rfloor$ . The nodes  $v_i$  when  $1 \leq i \leq 2q$ , are painted with color 1 when  $i$  is odd and color 2 when  $i$  is even and the nodes  $v_{i+\left\lfloor \frac{n}{2} \right\rfloor}$  when  $1 \leq i \leq 2q$ , are painted with color 3 when  $i$  is odd and color 4 when  $i$  is even. The node  $v_n$  is given color 5.

The nodes  $v_i$  when  $1 \leq i \leq 2q$ , dominate color class 2 when  $i$  is odd and dominate color class 1 when  $i$  is even. The nodes  $v_{i+\left\lfloor \frac{n}{2} \right\rfloor}$  when  $1 \leq i \leq 2q$ , dominate color class 4 when  $i$  is odd and color class 3 when  $i$  is even. The node  $v_n$  dominates color class 5.

Every adjacent node is given different color and also it is observed that every node of the graph dominates all the nodes of atleast one color class. Thus it is a dominator coloring of nodes and the dominator coloring number of  $\left\lfloor \frac{n}{2} \right\rfloor + 1$  regular graph on  $n$ (odd) nodes,  $n > 3$  and  $n \pmod 4 \equiv 1$  is given by  $\chi_d = 5$ .



**Figure 6:** Dominator coloring number of 6 - regular graph on 9 vertices is  $\chi_d(G) = 5$ .

$$\left( \text{Here } n = 9, n \pmod 4 \equiv 1 \text{ and } \left\lfloor \frac{n}{2} \right\rfloor + 1 = 6 \right)$$

### 3. CONCLUSION

In the field of graph theory, node coloring and domination have plenty of application in computer and communication network such as cell-phone network, mobile ad-hoc network, vehicular adhoc network, wireless sensor network etc. for routing and broadcasting the information between the nodes to the mobile devices. So the amalgamation of domination and coloring of graphs called the dominator coloring shall also be used in the study of interconnection networks; in particular to mobile ad-hoc network and vehicular adhoc network. In this paper, the dominator coloring number of some regular graphs are obtained.

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